

2-Dimensional Rigidity with Three Coincident Points

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joint work with Bill Jackson

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- 3 Results and Examples

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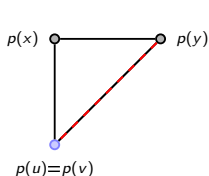
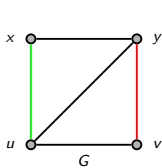
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- Can we characterise $\mathcal{R}_{uvw}(G)$ in a combinatorial way?

Two Coincident Points Example

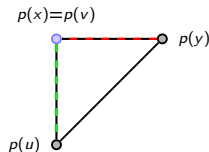
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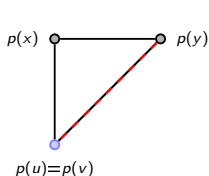
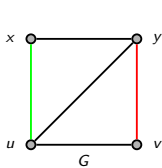
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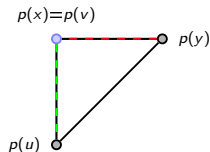
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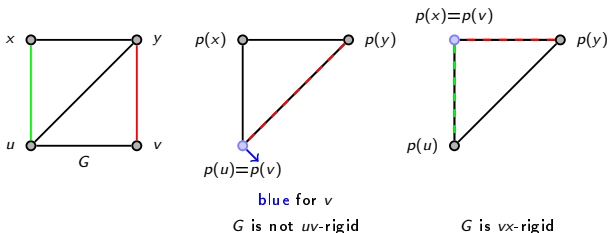
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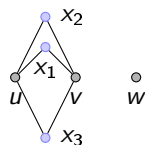
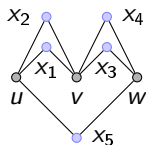
$$\begin{aligned} val_S(\mathcal{H}) &= \sum_{i=1}^k (2|H_i \setminus S| - 1) + 2(|S| - 1) \\ &= \sum_{i=1}^k (2|H_i| - 3) - 2(|S| - 1)(k - 1) \end{aligned}$$

S -Sparsity

- We say that G is S -sparse if, for all $H \subseteq V$ with $|H| \geq 2$, we have $i_G(H) \leq \text{val}_S(H)$, and for all S -compatible families \mathcal{H} , we have $i_G(\mathcal{H}) \leq \text{val}_S(\mathcal{H})$. It follows that, if G is S -sparse, then there is no edge between any distinct pair of vertices in S as $\text{val}_S(H) = 0$ for a set $H \subseteq S$.

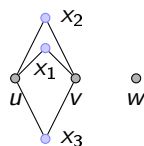
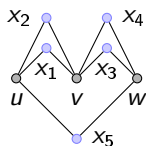
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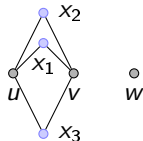
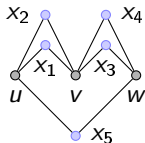
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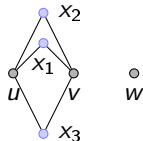
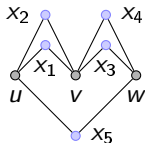
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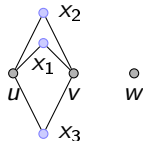
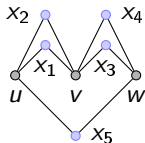
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 - $|X_i \cap X_j| \leq 1$ for all pairs $1 \leq i < j \leq k$.The collection $\mathcal{L} = \{\mathcal{H}, X_1, \dots, X_k\}$ is *thin* if (i) holds and
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 - $|X_i \cap \bigcup_{j=1}^t H_j| \leq 1$ for all $1 \leq i \leq k$.
- We define the value of \mathcal{L} as

$$\text{val}(\mathcal{L}) = \text{val}_S(\mathcal{H}) + \sum_{i=1}^k 2|X_i| - 3.$$

Theorem 1 (Fekete, Jordán and Kaszanitzky ($|U| = 2$) / Jackson, G. ($|U| \geq 3$))

Let $G = (V, E)$ be a graph and $U \subseteq V$. Then the family $\mathcal{I}_G := \{F \subseteq E, H = (V, F) \text{ is } (U) - \text{sparse}\}$ is a family of independent sets of a matroid, $\mathcal{M}_U(G)$ on E . Moreover, the rank of a set $E' \subseteq E$ in $\mathcal{M}_U(G)$ is equal to

$$\min\{\text{val}(\mathcal{L}) : \mathcal{L} \text{ is a thin cover of } E'\}.$$

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Let $G = (V, E)$ be a graph and $u, v, w \subseteq V$ be distinct. Then $\mathcal{M}_{uvw}(G) = \mathcal{R}_{uvw}(G)$.

Theorem 4 (Fekete, Jordán and Kaszanitzky)

Let $G = (V, E)$ be a graph and $u, v \in V$ be distinct vertices. Then G is uv -rigid if and only if $G - uv$ and G_{uv} are both rigid.

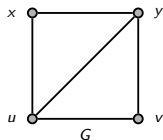
The graph G_{uv} is obtained from G by contracting u and v , and deleting double edges if there are any.

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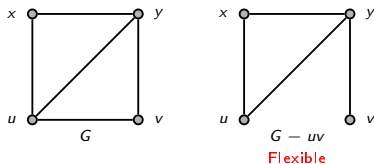


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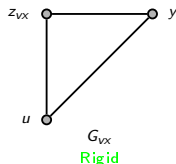
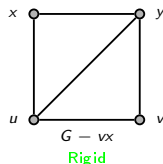
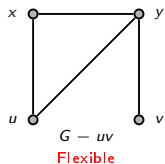
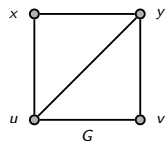


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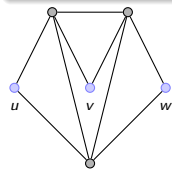
*Let $G = (V, E)$ be a graph and let $u, v, w \in V$ be distinct vertices and $G' = G - uv - uw - vw$. Then G is *uvw-rigid* if and only if G' and G'_S are rigid for all $S \subseteq \{u, v, w\}$ with $|S| \geq 2$.*

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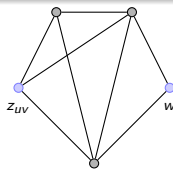
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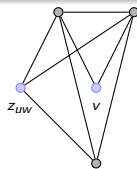
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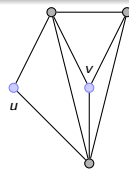
$G = G'$
Rigid



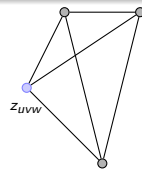
G_{uv}
Rigid



G_{uw}
Rigid



G_{vw}
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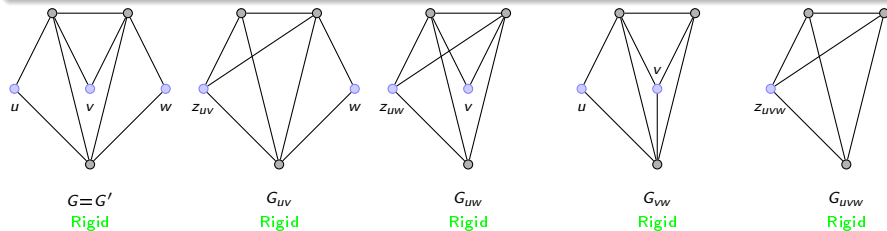
G_{uvw}
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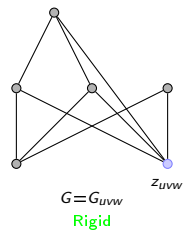
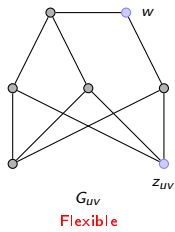
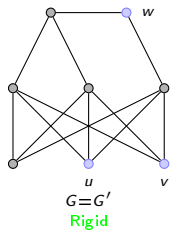
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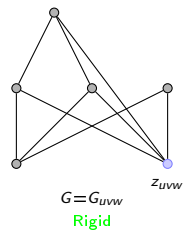
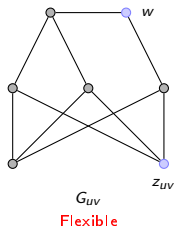
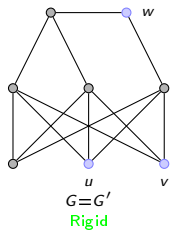


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- G is not uvw -rigid.

Thank you!